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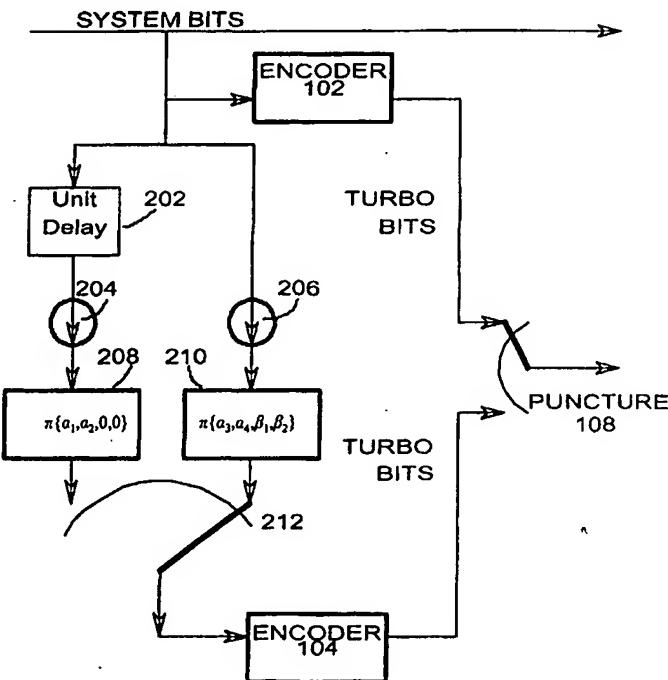
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(54) Title: INTERLEAVER USING CO-SET PARTITIONING

(57) Abstract

Construction of a Turbo code interleaver for blocks of any size is facilitated by partitioning the block into co-sets of predetermined size, permuting each co-set with an interleaver whose parameters are selected and optimized for the predetermined size, and recombining the elements of the permuted co-sets according to a predetermined order. Partitioning of the block into co-sets is accomplished by choosing a value n for the number of co-sets, and assigning to each co-set those positions of the block having a common value for modulo-n of the block size. The permuted output may be punctured to maintain a desired ratio.



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INTERLEAVER USING CO-SET PARTITIONING

FIELD OF THE INVENTION

This invention relates to communications systems and, more particularly, to interleavers for performing code modulation.

BACKGROUND OF THE INVENTION

10 Encoding communication channels, i.e., coded modulation, has been found to improve the bit error rate (BER) of electronic communication systems such as modem and wireless communication systems. Turbo coded modulation has proven to be a practical, power-efficient, and bandwidth-efficient modulation method for "random-error" channels characterized by additive white Gaussian noise (AWGN) or fading.

15 These random-error channels can be found, for example, in the Code Division Multiple Access (CDMA) environment.

The key innovation of Turbo Codes is an interleaver which permutes the original data frame before input to a second encoder. A conventional parallel concatenated Turbo Encoder is shown in Fig. 1. The Turbo Encoder includes two 20 constituent encoders 102 and 104, and one interleaver 106. Switch 108 can puncture alternate bits from each of encoders 102 and 104 so as to maintain the desired ratio (e.g., $R = \frac{1}{2}$). The construction and operation of the constituent encoder and the selection of a ratio are well known and thus will not be discussed further herein.

25 However the construction of the interleaver remains an open issue; especially for the small block size of Turbo Code.

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Presently, a conventional "random interleaver" is considered the best if the data frame size goes to infinity in an Additive White Gaussian Noise(AWGN) channel such as a CDMA channel. However, the best interleaver for data frames of finite size (i.e. Turbo Code) is yet to be determined.

5 Accordingly, there exists a need for interleaver designs for data frames of finite size. It is thus an advantage of the present invention to provide interleavers for data frames of finite size (i.e. Turbo Code). It is another advantage of the present invention to provide optimized interleavers for frames of predetermined size.

SUMMARY OF THE INVENTION

10 An aspect of the invention provides a method of permuting a block containing N data elements, where N is a positive integer greater than 1 and integrally divisible by a positive integer n greater than 1. The method includes partitioning the block into n co-sets each of size N/n ; permuting each co-set; and forming a permuted block by combining elements of the permuted co-sets in a predetermined order.

15 An aspect of the invention provides a system and method for partitioning said block into n co-sets by assigning to each of the n co-sets elements whose relative positions in said block have a common value of (N modulo(n)).

An aspect of the invention provides for permuting each co-set in an algebraic interleaver building block adapted to the size of the co-set and for applying co-set partitioning and interleaving several times. An aspect of the invention provides for permuting each co-set by generating an m-sequence according to the co-set.

Another aspect of the invention provides a system and method of producing Turbo code from a stream of N system bits, where N is a positive integer greater than

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1 and integrally divisible by a positive integer n greater than 1. The method includes
encoding the N system bits in a first encoder to produce a first stream of N Turbo bits;
partitioning the stream of N system bits into n co-sets each of size N/n ; permuting
each co-set; forming a permuted stream of N system bits by combining permuted said
5 co-sets in a predetermined order; encoding the N permuted system bits in a second
encoder to produce a second stream of N Turbo bits; and combining the first and
second streams of N Turbo bits in a predetermined order.

An aspect of the invention provides for partitioning the stream of N system
bits into n co-sets by assigning to each of said n co-sets elements whose relative
10 positions in the stream have a common value of (N modulo(n)) and for applying the
co-set partitioning and interleaving several times.

An aspect of the invention provides for producing two streams of Turbo bits
from a stream of system bits, and puncturing the two streams of N Turbo bits to
maintain a predetermined ratio of system bits to Turbo bits.

15 The invention will next be described in connection with certain exemplary
embodiments; however, it should be clear to those skilled in the art that various
modifications, additions, and subtractions can be made without departing from the
spirit or scope of the claims.

BRIEF DESCRIPTION OF THE DRAWINGS

20 The invention will be more clearly understood by reference to the following
detailed description of an exemplary embodiment in conjunction with the
accompanying drawings, in which:

Fig. 1 is a block diagram of a conventional parallel concatenated Turbo

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Encoder;

Fig. 2 is a block diagram of a parallel concatenated Turbo Encoder using algebraic co-set interleavers according to the present invention;

Fig. 3 depicts data loss due to puncturing in the conventional Turbo Encoder
5 of Fig. 1;

Fig. 4 is a block diagram of a parallel concatenated Turbo Encoder using m-
sequence generators as co-set interleavers according to the present invention.

DETAILED DESCRIPTION OF THE INVENTION

The present invention discloses an interleaver that improves the performance
10 of conventional Turbo Code. An embodiment of the invention employs an algorithm
to construct an arbitrary size algebraic interleaver based on smaller size interleavers
by exploiting the concept of mapping by set partitioning. This technique allows for
the parameterization of the algebraic interleaver and for interleaver optimization based
on selected criteria.

15 The described embodiment compensates for the loss of parity bits caused by
puncturing of a parallel concatenated Turbo Code. By employing a co-set partitioned
algebraic interleaver, the performance of punctured Turbo Code (especially for $R=1/2$)
is improved. The invention may be realized in software or hardware (i.e. application
specific integrated circuits (ASIC), programmable logic arrays (PLA), or any other
20 suitable logic device).

While the present disclosure discusses an embodiment of interleavers with
respect to Turbo Codes and cdma2000, those skilled in the art will appreciate that
these interleavers may also be used in other systems such as present-day CDMA, and

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with other codes as well.

Construction of Algebraic Interleaver Based on Co-Set Partitioning

The design of the present interleaver is based on a deterministic algebraic structure. It employs the idea of co-set partitioning to design an algebraic interleaver 5 that can interleave blocks of arbitrary size by partitioning the block into smaller co-set sub-blocks and using smaller sized algebraic interleavers to interleave the co-sets. Elements from the interleaved co-sets are then combined in a predetermined order to form the interleaved version of the original block.

The co-set partitioning is accomplished in a present embodiment by assigning 10 to each of n co-sets elements of the block that share the same residual of a modulo-n calculation.

For example, a block containing 12 elements identified as

{1 2 3 4 5 6 7 8 9 10 11 12}

could be partitioned into 2 co-sets containing
15 {1 3 5 7 9 11} (modulo-2 residual = 1)
 {2 4 6 8 10 12} (modulo-2 residual = 0)

or into 3 co-sets containing
20 {1 4 7 10} (modulo-3 residual = 1)
 {2 5 8 11} (modulo-3 residual = 2)
 {3 6 9 12} (modulo-3 residual = 0)

or into 4 co-sets containing
25 {1 5 9} (modulo-4 residual = 1)
 {2 6 10} (modulo-4 residual = 2)
 {3 7 11} (modulo-4 residual = 3)
 {4 8 12} (modulo-4 residual = 0).

(Although the block could also be partitioned into 6 co-sets of two elements each, the 30 results of permuting the two-element co-sets would probably not be of practical use.)

An example will now be considered in detail of interleaving a block having 48

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data elements, identified as elements 1 through 48. The block will be divided into two co-sets, one containing those of elements 1 through 48 that have a modulo-2 residue of 0 (i.e., the 24 "even-numbered" elements 2, 4, 6, ... 48) and one containing those that have a modulo-2 residue of 1 (i.e., the 24 "odd-numbered" elements 1, 3, 5, ... 47). The two blocks may be referred to in this discussion as the even block and the odd block. Those skilled in the art will appreciate that the choice of two as the number of co-sets is exemplary and not limiting. Each of the two blocks will be permuted row-by-row and column-by-column.

Each of the 24-element blocks is treated as a four-row by six-column array--
10 i.e., as an $N_1 \times N_2$ array where N_1 is chosen to be 4 and N_2 is accordingly chosen to be 6. Those skilled in the art realize that these numbers are exemplary and that other numbers could be chosen. Prime numbers P_1 and P_2 are then chosen. P_1 is a prime number greater than N_1 ; it is preferably the least prime number greater than N_1 , but the method will work if greater prime numbers are chosen. Accordingly, a value of 5 is
15 chosen for P_1 , since 5 is the least prime number greater than 4 (which is the value of N_1). Similarly, a value of 7 is chosen for P_2 , since 7 is the least prime number greater than 6 (which is the value of N_2). Values of α_1 , α_2 , β_1 and β_2 are then chosen: values α_1 and α_2 are initial roots of P_1 and P_2 , respectively, found by conventional methods. β_1 and β_2 may be selected for optimum results from

20 $\beta_1 \leq P_1-1$ and $\beta_2 \leq P_2-1$.

In the present example, $\alpha_1 = 2$, $\alpha_2 = 3$, $\beta_1 = 1$, and $\beta_2 = 3$.

The block to be interleaved is read in to the N_1 by N_2 array in a row-by-row manner. Thus, the even block is read in as

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02	04	06	08	10	12
14	16	18	20	22	24
26	28	30	32	34	36
38	40	42	44	46	48

5

Then the row indices $i_{row}(n)$ are permuted according to:

for $n = 1, 2, \dots, P_1-1$
 10 $i_{row}(n) = \alpha_1^{n+\beta_1} \bmod (P_1)$
 puncture $i_{row}(n) > N_1$

Accordingly, for rows 1, 2, 3, and 4 of the present example:

15 $i_{row}(1) = 2^{1+1} \bmod 5 = 4 \bmod 5 = 4$
 $i_{row}(2) = 2^{2+1} \bmod 5 = 8 \bmod 5 = 3$
 $i_{row}(3) = 2^{3+1} \bmod 5 = 16 \bmod 5 = 1$
 $i_{row}(4) = 2^{4+1} \bmod 5 = 32 \bmod 5 = 2$

20 Thus, the former row 1 becomes row 4, the former row 2 becomes row 3, etc.,
 yielding:

25 $26 \ 28 \ 30 \ 32 \ 34 \ 36$
 $38 \ 40 \ 42 \ 44 \ 46 \ 48$
 $14 \ 16 \ 18 \ 20 \ 22 \ 24$
 $02 \ 04 \ 06 \ 08 \ 10 \ 12$

Next, the column indices $i_{col}(n)$ are permuted according to:

30 for $n = 1, 2, \dots, P_2-1$
 $i_{col}(n) = \alpha_2^{n+\beta_2} \bmod (P_2)$
 puncture $i_{col}(n) > N_2$

Accordingly, for columns 1, 2, 3, 4, 5, and 6 of the present example:

35 $i_{col}(1) = 3^{1+3} \bmod 7 = 81 \bmod 7 = 4$
 $i_{col}(2) = 3^{2+3} \bmod 7 = 243 \bmod 7 = 5$
 $i_{col}(3) = 3^{3+3} \bmod 7 = 729 \bmod 7 = 1$
 $i_{col}(4) = 3^{4+3} \bmod 7 = 2187 \bmod 7 = 3$
 $i_{col}(5) = 3^{5+3} \bmod 7 = 6561 \bmod 7 = 2$
 40 $i_{col}(6) = 3^{6+3} \bmod 7 = 19683 \bmod 7 = 6$

Thus, the former column 1 becomes column 4, the former column 2 becomes column 5, etc., yielding:

45 $30 \ 34 \ 32 \ 26 \ 28 \ 36$
 $42 \ 46 \ 44 \ 38 \ 40 \ 48$
 $18 \ 22 \ 20 \ 14 \ 16 \ 24$
 $06 \ 10 \ 08 \ 02 \ 04 \ 12$

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5 Next, a cyclic shift of the rows by β_1 rows (1 row) is performed, yielding:

10

42	46	44	38	40	48
18	22	20	14	16	24
06	10	08	02	04	12
30	34	32	26	28	36

Next, a cyclic shift of the columns by β_2 columns (3 columns) is performed, yielding:

15

38	40	48	42	46	44
14	16	24	18	22	20
02	04	12	06	10	08
26	28	36	30	34	32

20 Reading out the interleaved even block in a column-by-column manner results in the interleaver output index of:

38, 14, 02, 26, 40, 16, 04, 28, 48, 24, 12, 36, 42, 18, 06, 30, 46, 22, 10, 34, 44, 20, 08, 32.

Similar processing is employed to interleave the odd block (the 24 odd elements of the 48-element block being considered). The interleaved co-sets are recombined according to a predetermined order. Blocks of a size that is not numerically convenient to handle may be padded to a convenient size for permuting, and the padding positions may be discarded after permuting. For example, a block of 45 elements could have extra elements 46, 47, and 48 added before being permuted as in the present example, and elements 46, 47, and 48 could be discarded after permuting.

25 Thus, a basic algebraic interleaver building block of a particular size (here, 24) has been given parameters. The parameters may be expressed as $\pi\{\alpha_1, \alpha_2, \beta_1, \beta_2\}$, where π connotes the interleaver block, and the roots $\{\alpha_1, \alpha_2\}$ and shifts $\{\beta_1, \beta_2\}$ can

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be selected against different optimization criteria specified for an algebraic interleaver and refined to their optimum values for the building block size.

An example of a Turbo Encoder using such a co-set interleaver is shown in Fig. 2. In this example the input sequence is split into two co-sets, one comprising the 5 even-numbered bits of the input sequence and one comprising the odd-numbered bits. Elements 204 and 206 are decimators that pass every second element routed to them. In conjunction with unit delay 202 in the path to only one of decimators 204 and 206, this has the effect of passing a co-set consisting of the sequence of odd system bits (1, 3, ...) to interleaver building block 210, and a co-set consisting of the sequence of 10 even system bits to interleaver building block 208. Those skilled in the art realize that other passing schemes can be used.

The starting positions of switches 108 and 212 are initialized to the positions shown in Fig. 2. They then toggle synchronously with the flow of bits through the Turbo Encoder of Fig. 2. Switch 212 thus reconstitutes a single block for input to 15 encoder 104, consisting of the first element output from interleaver 210, the first element output from interleaver 208, the second element output from interleaver 210, and so on. In the absence of switch 108 the configuration of Fig. 2 would output a bit stream with a ratio $R = 1/3$ (one system bit and two Turbo bits out for every system bit in). Switch 108 punctures alternate Turbo bits from each of encoders 102 and 104 20 so that the output has a ratio $R = 1/2$.

The algebraic interleaver building block (co-set) is chosen with size $N_0 = N/2$, where N is the size of the block to be interleaved. The choice of $N/2$ for the value of N_0 is a design choice; those skilled in art realize that other building-block sizes are

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possible. Each building block is then assigned a number N_1 of rows and a number N_2 of columns, with $N_1 \times N_2 = N_0$. For the first building block the parameters chosen are $\{\alpha_1, \alpha_2, 0, 0\}$, and for the second building block the chosen parameters are $\{\alpha_3, \alpha_4, \beta_1, \beta_2\}$. The parameterized interleaver can be optimized with appropriate values with respect to the Turbo Code performance. Table 1 illustrates a set of optimized parameters determined for the Turbo Codes specified in cdma2000.

Data Rate (kbps)	Frame Size N	N_1, N_2	P_1, P_2	$\alpha_1, \alpha_2, \alpha_3, \alpha_4, \beta_1, \beta_2$
9.6	192	16,6	17,7	7,5,5,3,7,2
14.4	288	16,9	17,11	6,6,3,6,2,4
19.2	384	16,12	17,13	3,6,6,6,14,6
28.8	576	16,18	17,19	3,13,6,13,0,9
38.4	768	16,24	17,29	3,2,6,2,0,11
57.6	1152	16,36	17,37	2,3,20,3,29,6

Table 1

In alternative practices of the invention, co-set partitioning may be performed more than once. In the above example, where the columns 1, 2, 3, 4, 5, and 6 of the even co-set are permuted to columns 4, 5, 1, 3, 2, and 6, further co-set partitioning might be performed. That is, taking the new first, third, and fifth columns as one co-set and the new second, fourth, and sixth columns as another co-set, thereby producing two co-sets,

15	30 32 28	34 26 36
	42 44 40	46 38 48
	18 20 16	22 14 24
	06 08 04	10 02 12

20 which are then each subjected to row and column cyclic shifts as in the above example. In other practices of the invention, each row or column shift may use a

different value. That is, the value of β can be row-independent, or column-independent, or both.

Punctured Turbo Codes With No Loss of Turbo Code Bits

5 An additional advantage gained from practicing the present invention's co-set partitioning with the number of co-sets chosen as two is that puncturing to attain a particular ratio R (as by switch 108 of Fig. 2) will not puncture both Turbo Code bits corresponding to a particular system bit. Fig. 3 shows that for the conventional Turbo Encoder of Fig. 1, showing five system bits connoted S_1 through S_5 , and taking 10 system bit S_2 as an example, its corresponding Turbo Code bits T_{12} and T_{22} are both lost in puncturing (for an assumed order of interleaved bits). Since one of those bit streams is being interleaved relative to the other, it cannot be guaranteed that this will not occur. When decoding at a destination, there is a decreased likelihood of correct interpretation of system bit 2.

15 According to the present invention wherein the input sequence is partitioned into an even sequence and an odd sequence, interleaving is performed separately on each sequence, and the sequences are subsequently recombined, each system bit will retain one of its corresponding Turbo Code bits after puncturing.

Further, a single pseudo random sequence generator (e.g., m-sequence, M-sequence, Gold sequence, Kasami sequence, etc.) can be employed as the interleaver building block, or as illustrated in Fig. 4, two m-sequence generators 408 and 410 are combined to form an interleaver for co-set partitioning.

It will thus be seen that the invention provides advantageous interleaving for use in a Turbo Encoder for encoding frames of finite size. Those skilled in the art will

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appreciate that the configurations depicted in Figs. 2 and 4 provide Turbo Encoders with co-set partitioning.

It will be understood that changes may be made in the above construction and in the foregoing sequences of operation without departing from the scope of the invention. It is accordingly intended that all matter contained in the above description or shown in the accompanying drawings be interpreted as illustrative rather than in a limiting sense.

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WHAT IS CLAIMED IS:

1. A method of permuting a block of data elements, comprising:
partitioning said block into a plurality of co-sets;
permuting each co-set; and
5 forming a permuted block by combining elements of said permuted co-sets in a
predetermined order.

2. The method according to Claim 1 wherein:
said block contains N data elements, where N is a positive integer greater than 1 and
10 integrally divisible by n, a positive integer greater than 1; and
the step of partitioning said block into co-sets comprises assigning to each of n co-sets
elements whose relative positions in said block have a common value of (N
modulo(n)).

15 3. The method according to Claim 2 wherein the value of n is 2,
whereby a first co-set contains all odd-numbered elements of said block and a second
co-set contains all even-numbered elements of said block.

4. The method according to Claim 1 wherein the step of permuting each co-set
20 comprises:
reading each co-set in a row-by-row manner into an array having N_1 columns and N_2
rows where $N_1 \times N_2$ is equal to N/n where N is a number of elements in said block
and n is a number of co-sets;

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choosing prime numbers $P_1 < N_1$ and $P_2 < N_2$;

determining α_1 and α_2 as the roots of P_1 and P_2 ;

selecting values β_1 and β_2 ;

permuting the row indices of said array as a function of α_1 and β_1 ;

5 permuting the column indices of said array as a function of α_2 and β_2 ;

cyclically shifting the rows of said array by β_1 rows;

cyclically shifting the columns of said array by β_2 columns; and

reading each co-set out of said array in a column-by-column manner.

10 5. The method according to Claim 1 wherein the step of permuting each co-set comprises:

reading each co-set in a row-by-row manner into an array having N_1 columns and N_2 rows where $N_1 \times N_2$ is equal to N/n where N is a number of elements in said block and n is a number of co-sets;

15 choosing prime numbers $P_1 < N_1$ and $P_2 < N_2$;

determining α_1 and α_2 as the roots of P_1 and P_2 ;

selecting values β_1 and β_2 ;

permuting the row indices of said array as a function of α_1 and β_1 ;

permuting the column indices of said array as a function of α_2 and β_2 ;

20 cyclically shifting each row of said array by a predetermined number of rows;

cyclically shifting each column of said array by a predetermined number of columns;

and

reading each co-set out of said array in a column-by-column manner.

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6. The method according to Claim 1 wherein the step of permuting each co-set comprises generating an m-sequence according to each co-set.
7. A method of producing Turbo code from a stream of system elements, comprising;
 - 5 encoding the system elements in a first encoder to produce a first stream of Turbo elements;
 - partitioning said stream of system elements into a plurality of co-sets;
 - permuting each co-set;
 - forming a permuted stream of system elements by combining said permuted co-sets in
 - 10 a predetermined order;
 - encoding the permuted system elements in a second encoder to produce a second stream of Turbo elements; and
 - combining said first and second streams of Turbo elements in a predetermined order.
- 15 8. The method according to Claim 7 wherein the step of partitioning said stream of system elements into co-sets comprises assigning to each of n co-sets elements whose relative positions in said stream have a common value of (N modulo(n)) where N is a number of elements in said stream of system elements.
- 20 9. The method according to Claim 7 wherein said combining of said first and second streams of N turbo elements includes puncturing of said first and second streams of Turbo elements to maintain a predetermined ratio of system elements to Turbo elements.

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10. The method according to Claim 9 wherein said predetermined ratio is $\frac{1}{2}$ and the value of n is 2, whereby said Turbo code includes at least one element corresponding to each system element.

5

11. Apparatus for permuting a block of data elements, comprising:
means for partitioning said block into a plurality of co-sets;
means for permuting each co-set; and
means for forming a permuted block by combining elements of said permuted co-sets
10 in a predetermined order.

12. The apparatus according to Claim 11 wherein:
said block contains N data elements, where N is a positive integer greater than 1 and
integrally divisible by n, a positive integer greater than 1; and
15 the means for partitioning said block into co-sets assigns to each of n co-sets elements
whose relative positions in said block have a common value of (N modulo(n)).

13. The apparatus according to Claim 12 wherein the value of n is 2,
whereby a first co-set contains all odd-numbered elements of said block and a second
20 co-set contains all even-numbered elements of said block.

14. The apparatus according to Claim 11 wherein the means for permuting each co-set comprises:

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means for reading each co-set in a row-by-row manner into an array having N_1 columns and N_2 rows where $N_1 \times N_2$ is equal to N/n where N is a number of elements in said block and n is a number of co-sets;

means for choosing prime numbers $P_1 < N_1$ and $P_2 < N_2$;

5 means for determining α_1 and α_2 as the roots of P_1 and P_2 ;

means for selecting values β_1 and β_2 ;

means for permuting the row indices of said array as a function of α_1 and β_1 ;

means for permuting the column indices of said array as a function of α_2 and β_2 ;

means for cyclically shifting the rows of said array by β_1 rows;

10 means for cyclically shifting the columns of said array by β_2 columns; and

means for reading each co-set out of said array in a column-by-column manner.

15. The apparatus according to Claim 11 wherein the means for permuting each co-set comprises:

15 means for reading each co-set in a row-by-row manner into an array having N_1 columns and N_2 rows where $N_1 \times N_2$ is equal to N/n where N is a number of elements in said block and n is a number of co-sets;

means for choosing prime numbers $P_1 < N_1$ and $P_2 < N_2$;

means for determining α_1 and α_2 as the roots of P_1 and P_2 ;

20 means for selecting values β_1 and β_2 ;

means for permuting the row indices of said array as a function of α_1 and β_1 ;

means for permuting the column indices of said array as a function of α_2 and β_2 ;

means for cyclically each row of said array by a predetermined number of rows;

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means for cyclically shifting each column of said array by a predetermined number of columns; and

means for reading each co-set out of said array in a column-by-column manner.

5 16. The apparatus according to Claim 11 wherein the means for permuting each co-set comprises means for generating an m-sequence according to each co-set.

17. Apparatus for producing Turbo code from a stream of system elements, comprising;

10 means for encoding the system elements to produce a first stream of Turbo elements; means for partitioning said stream of system elements into a plurality of co-sets; means for permuting each co-set; means for forming a permuted stream of system elements by combining said permuted co-sets in a predetermined order;

15 means for encoding the permuted system elements to produce a second stream of Turbo elements; and

means for combining said first and second streams of Turbo elements in a predetermined order.

20 18. The apparatus according to Claim 17 wherein the means for partitioning said stream of system elements into co-sets assigns to each of n co-sets elements whose relative positions in said stream have a common value of (N modulo(n)) where N is a number of elements in said stream of system elements.

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19. The apparatus according to Claim 17 wherein said means for combining said first
and second streams of N turbo elements punctures said first and second streams of
Turbo elements to maintain a predetermined ratio of system elements to Turbo
5 elements.

20. The apparatus according to Claim 19 wherein said predetermined ratio is $\frac{1}{2}$ and
the value of n is 2, whereby said Turbo code includes at least one element
corresponding to each system element.

10

21. A system for permuting a block of data elements, comprising:
partitioning logic for partitioning said block into a plurality of co-sets;
permuting logic for permuting each co-set; and
combining logic for forming a permuted block by combining elements of said
15 permuted co-sets in a predetermined order.

22. The system according to Claim 21 wherein:
said block contains N data elements, where N is a positive integer greater than 1 and
integrally divisible by n, a positive integer greater than 1; and
20 said partitioning logic assigns to each of n co-sets elements whose relative positions in
said block have a common value of (N modulo(n)).

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23. The system according to Claim 22 wherein the value of n is 2,
whereby a first co-set contains all odd-numbered elements of said block and a second
co-set contains all even-numbered elements of said block.

5 24. The system according to Claim 21 wherein said permuting logic comprises:

read-in logic storing each co-set in a row-by-row manner into an array having N_1
columns and N_2 rows where $N_1 \times N_2$ is equal to N/n where N is a number of
elements in said block and n is a number of co-sets;
arithmetic logic configured to:

10 choose prime numbers $P_1 < N_1$ and $P_2 < N_2$;

determine α_1 and α_2 as the roots of P_1 and P_2 ; and

select values β_1 and β_2 ,

data manipulation logic configured to:

permute the row indices of said array as a function of α_1 and β_1 ;

15 permute the column indices of said array as a function of α_2 and β_2 ;

cyclically shift the rows of said array by β_1 rows; and

cyclically shift the columns of said array by β_2 columns, and

read-out logic reading each co-set out of said array in a column-by-column manner.

20 25. The system according to Claim 21 wherein said permuting logic comprises:

read-in logic storing each co-set in a row-by-row manner into an array having N_1
columns and N_2 rows where $N_1 \times N_2$ is equal to N/n where N is a number of
elements in said block and n is a number of co-sets;

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arithmetic logic configured to:

choose prime numbers $P_1 < N_1$ and $P_2 < N_2$;

determine α_1 and α_2 as the roots of P_1 and P_2 ; and

select values β_1 and β_2 ,

5 data manipulation logic configured to:

permute the row indices of said array as a function of α_1 and β_1 ;

permute the column indices of said array as a function of α_2 and β_2 ;

cyclically shift each row of said array by a predetermined number of rows; and

cyclically shift each column of said array by a predetermined number of columns,

10 and

read-out logic reading each co-set out of said array in a column-by-column manner.

26. The system according to Claim 21 wherein the permuting logic generates an m-

sequence according to each co-set.

15

27. A system for producing Turbo code from a stream of system elements,

comprising;

a first encoder configured to encode the system elements to produce a first stream of

Turbo elements;

20 partitioning logic configured to partition said stream of system elements into a

plurality of co-sets;

permuting logic configured to permute each co-set;

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first combining logic forming a permuted stream of system elements by combining
said permuted co-sets in a predetermined order;

a second encoder encoding the permuted system elements to produce a second stream
of Turbo elements; and

5 second combining logic combining said first and second streams of Turbo elements in
a predetermined order.

28. The system according to Claim 27 wherein the partitioning logic assigns to each
of n co-sets elements whose relative positions in said stream have a common value
10 of (N modulo(n)) where N is a number of elements in said stream of system
elements.

29. The system according to Claim 27 wherein said second combining logic
punctures said first and second streams of Turbo elements to maintain a
15 predetermined ratio of system elements to Turbo elements.

30. The system according to Claim 29 wherein said predetermined ratio is $\frac{1}{2}$ and the
value of n is 2, whereby said Turbo code includes at least one element
corresponding to each system element.

20

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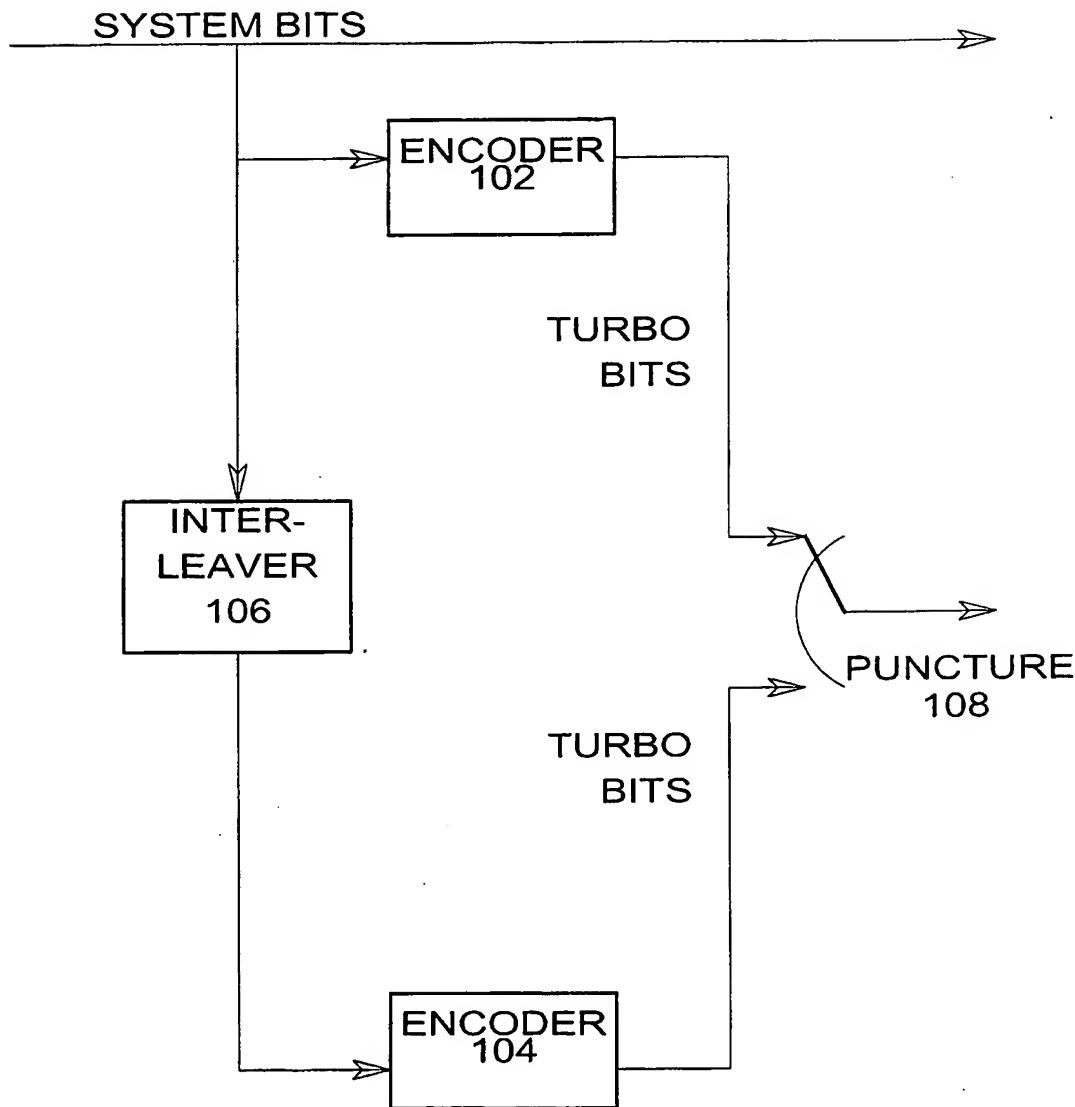


Fig. 1 (Prior Art)

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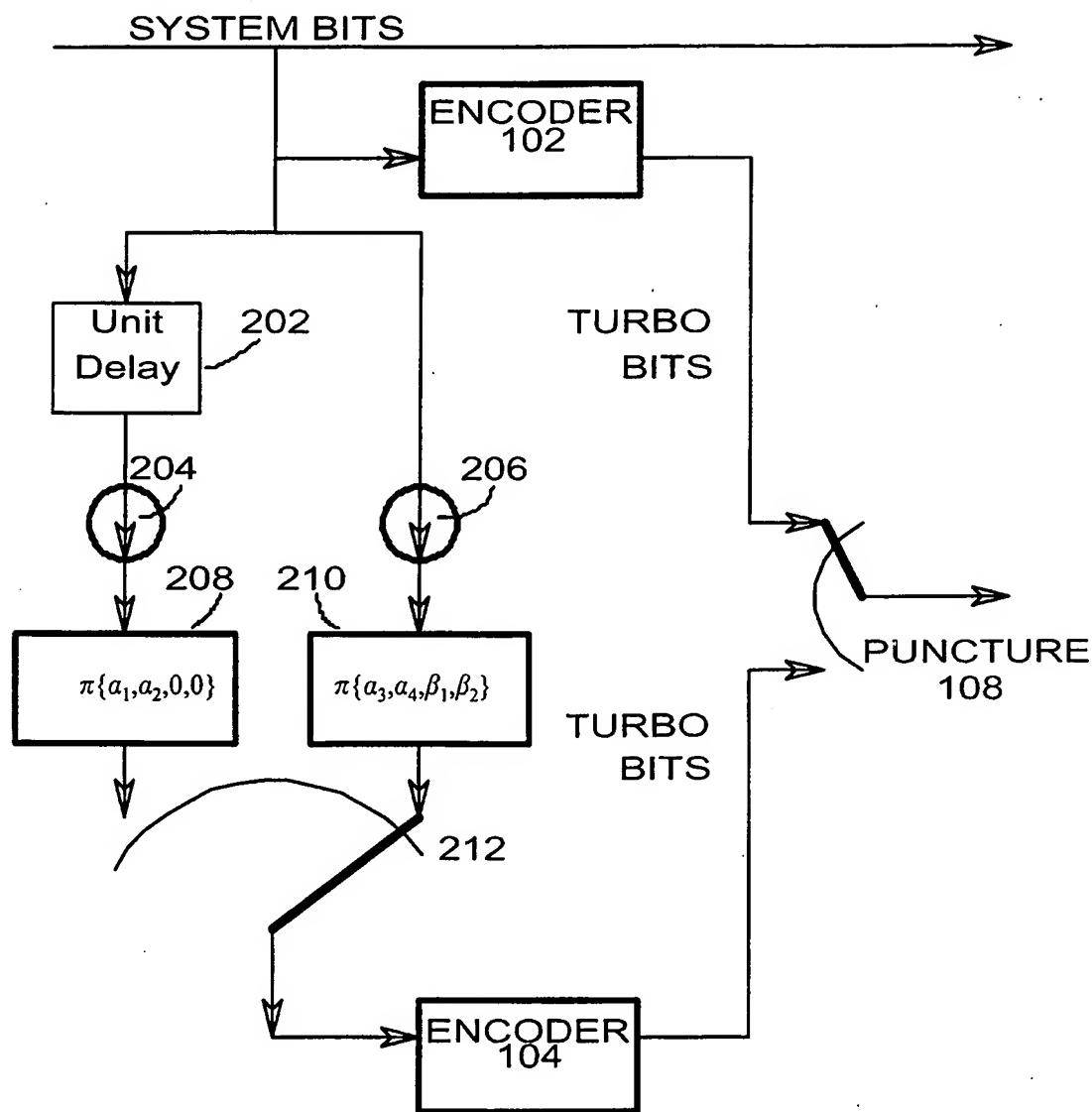


Fig. 2

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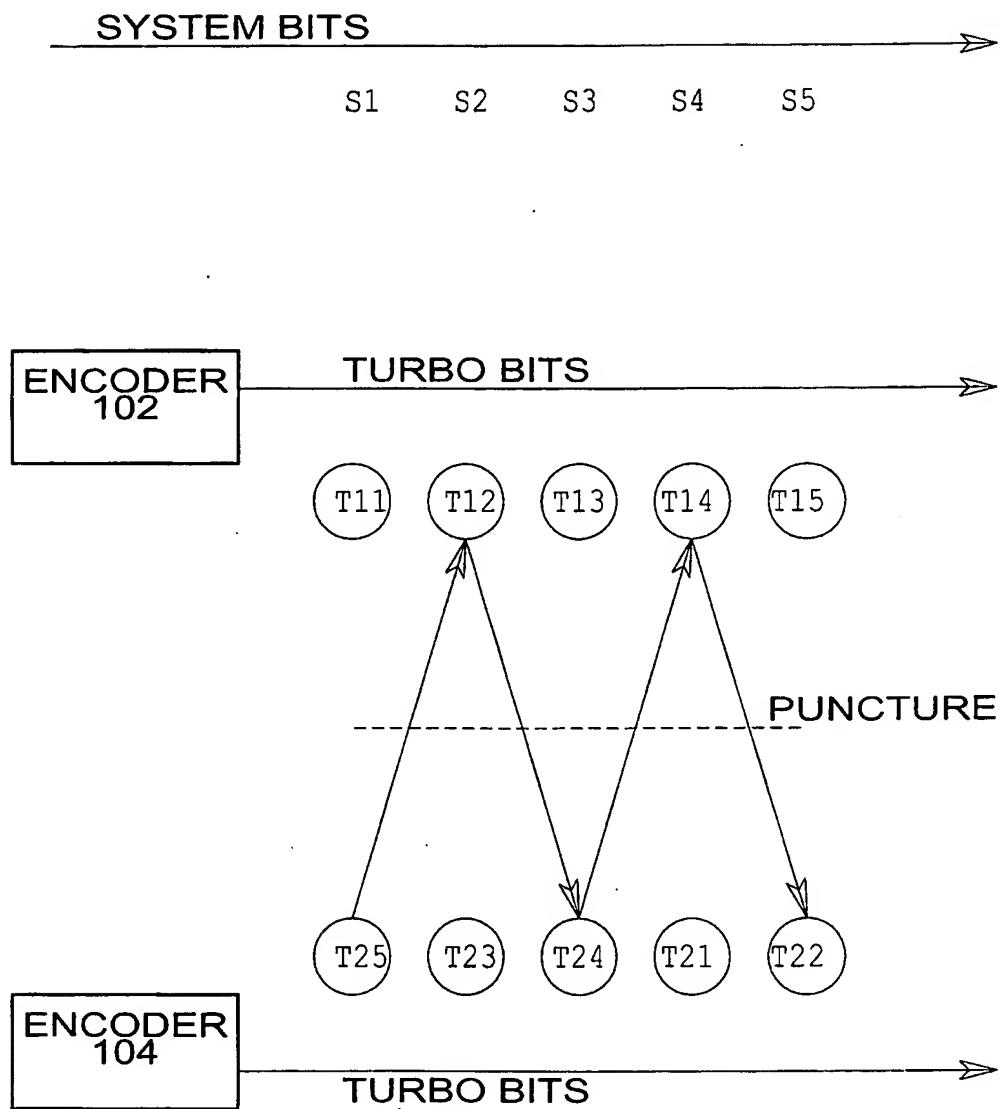


Fig. 3 (Prior Art)

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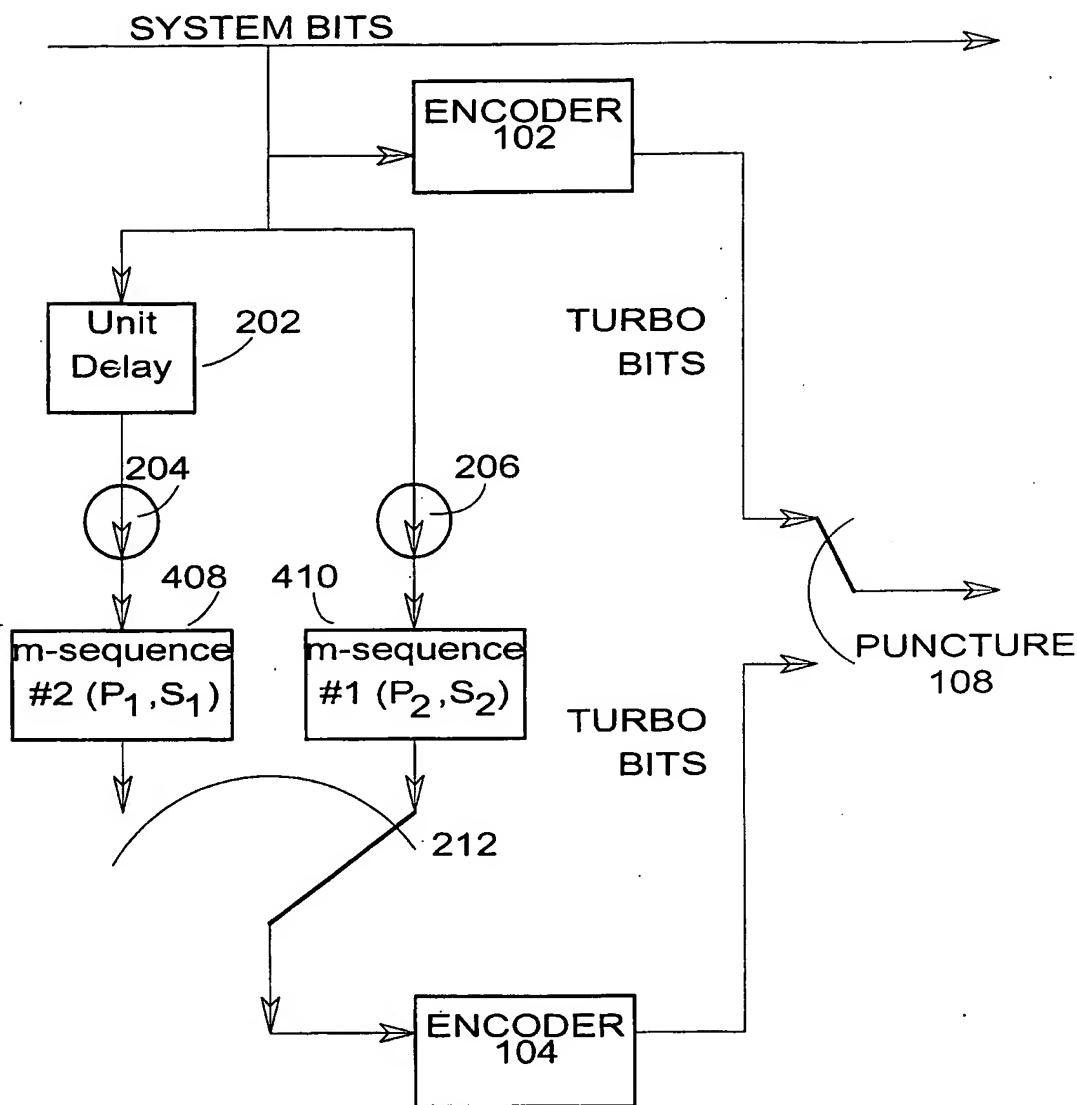


Fig. 4

INTERNATIONAL SEARCH REPORT

Int. Appl. No
PCT/IB 99/01594

A. CLASSIFICATION OF SUBJECT MATTER
IPC 7 H03M13/27 H03M13/29 H03M13/00

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 7 H03M

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	US 5 133 061 A (PFISTER GREGORY F ET AL) 21 July 1992 (1992-07-21) the whole document ---	1-30
P, Y	EP 0 928 071 A (CANON KK) 7 July 1999 (1999-07-07) the whole document ---	1-30
A	WO 97 40582 A (GEN ELECTRIC) 30 October 1997 (1997-10-30) the whole document -----	1-30

Further documents are listed in the continuation of box C.

Patent family members are listed in annex.

* Special categories of cited documents :

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Date of the actual completion of the international search 21 December 1999	Date of mailing of the international search report 12/01/2000
Name and mailing address of the ISA European Patent Office, P.B. 5818 Patentlaan 2 NL - 2280 HV Rijswijk Tel. (+31-70) 340-2040, Tx. 31 651 epo nl, Fax: (+31-70) 340-3016	Authorized officer Devergranne, C

INTERNATIONAL SEARCH REPORT

Information on patent family members

Intell. n al Application No

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